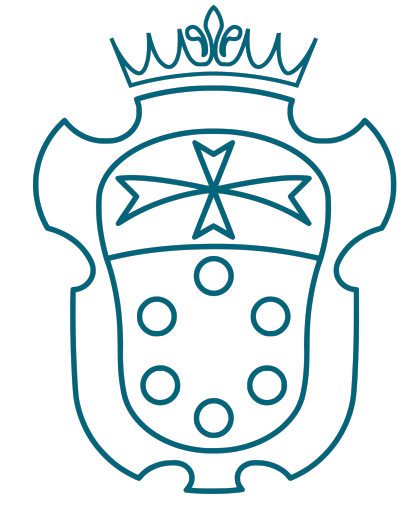




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**SCUOLA
NORMALE
SUPERIORE**

Random purification channel made simple

Filippo Girardi

Francesco Anna Mele

Ludovico Lami

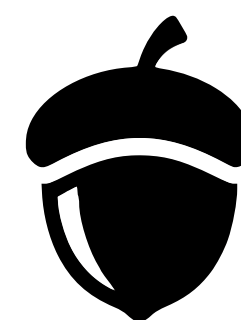
Scuola Normale Superiore, Pisa

Quantum Physics

[Submitted on 8 Oct 2025]

Conjugate queries can help

Ewin Tang, John Wright, Mark Zhandry



1.2 The acorn trick

Our main technique is what we call t
Suppose we have n copies of a p

arXiv > quant-ph > arXiv:2511.23451

Quantum Physics

[Submitted on 28 Nov 2025]

Random purification channel made simple

Filippo Girardi, Francesco Anna Mele, Ludovico Lami

arXiv > quant-ph > arXiv:2512.16878

arXiv > quant-ph > arXiv:2604.06325

Quantum Physics

[Submitted on 7 Apr 2026 (v1), last revised 7 May 2026 (this version, v2)]

Probabilistic and approximate universal quantum purification machines

Zoe G. del Toro, Jessica Bavaresco

arXiv > quant-ph > arXiv:2512.20599

arXiv > quant-ph > arXiv:2604.15269

Quantum Physics

[Submitted on 16 Apr 2026]

Cloning is as Hard as Learning for Stabilizer States

Nikhil Bansal, Matthias C. Caro, Gaurav Mahajan

arXiv > quant-ph > arXiv:2603.18136

Quantum Physics

[Submitted on 18 Mar 2026]

Towards sample-optimal learning of bosonic Gaussian quantum states

arXiv > quant-ph > arXiv:2605.21457

Quantum Physics

[Submitted on 20 May 2026]

An Exponential Sample-Complexity Advantage for Coherent Quantum Inference

Zhaoyi Li, Elias Theil, Aram W. Harrow, Isaac Chuang

arXiv >

arXiv > quant-ph > arXiv:2605.03975

Quantum Physics

arXiv > quant-ph > arXiv:2606.23592

Quantum Physics

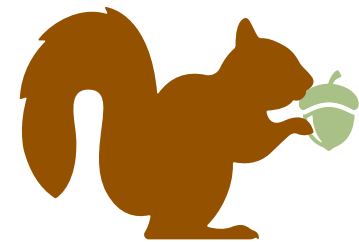
[Submitted on 22 Jun 2026]

Random dimension reduction and learning symmetric properties of quantum states

Angus Lowe, Xinyu Tan

Overview

$$\rho_A^{\otimes N} \xrightarrow{\quad \nearrow \quad} |\psi_\rho\rangle_{AE}^{\otimes N}$$



Why do we want to purify states?

(and why this is physically impossible)

How can we make this possible?

(random purification channel)

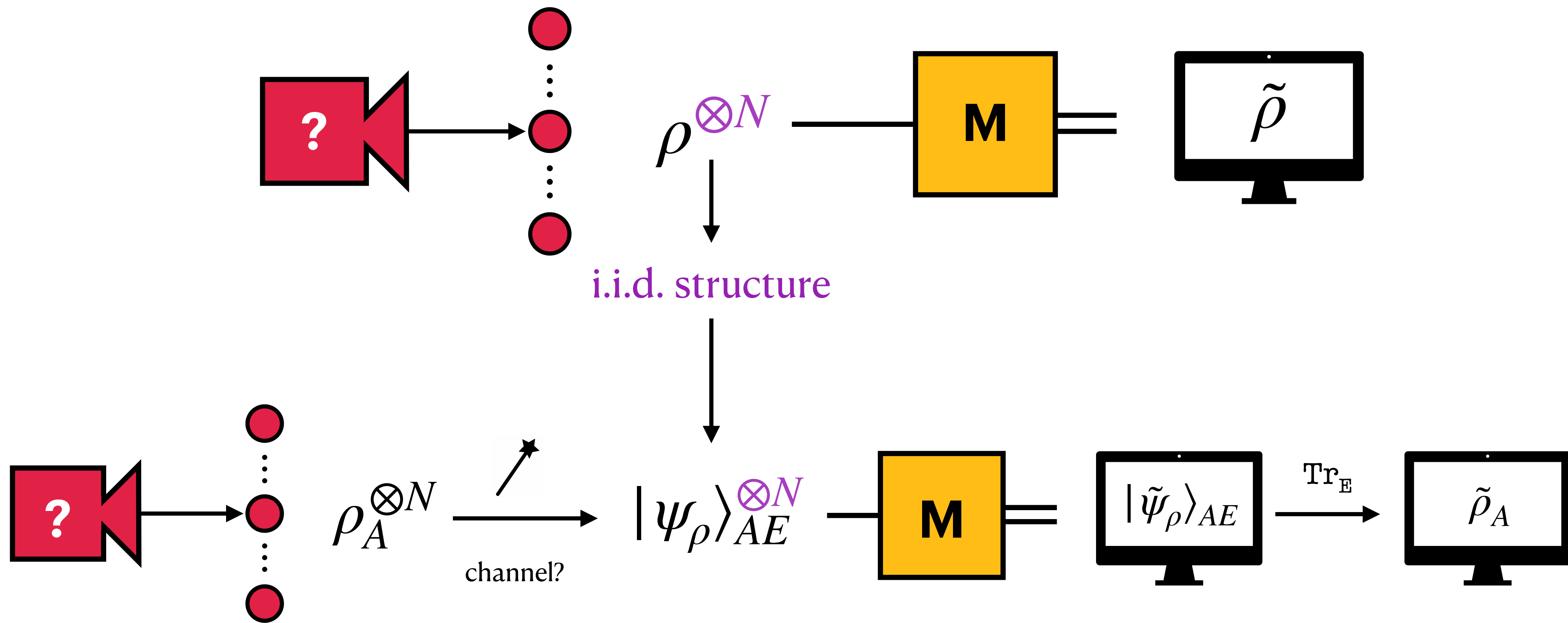
→ the simplest construction of the random purification channel

(without representation theory)

→ applications and generalisations

(new tool in quantum learning theory, but also in quantum Shannon theory)

Learning theory



No purifying channel

Physical implementation of a **universal purification channel** $\tilde{\Lambda}$?

$$\rho_A \xrightarrow{\tilde{\Lambda}} |\rho\rangle_{AB}$$

$$\sigma_A \xrightarrow{\tilde{\Lambda}} |\sigma\rangle_{AB}$$

$$\tilde{\Lambda}(\rho_A) = |\rho\rangle\langle\rho|_{AB}$$

$$\tilde{\Lambda}(\sigma_A) = |\sigma\rangle\langle\sigma|_{AB}$$

linearity

$$\tilde{\Lambda}\left(\frac{1}{2}\rho_A + \frac{1}{2}\sigma_A\right) = \frac{1}{2}|\rho\rangle\langle\rho|_{AB} + \frac{1}{2}|\sigma\rangle\langle\sigma|_{AB} \quad \text{mixed} \quad (\text{unless } |\rho\rangle = |\sigma\rangle)$$

Z. Liu, Z. Du, Z. Cai, and Z.-W. Liu. *No universal purification in quantum mechanics*. Preprint arXiv:2509.21111 (2025)

K. Fang and Z.-W. Liu. *No-go theorems for quantum resource purification*. Phys. Rev. Lett., 125(6):060405, (2020)

K. Fang and Z.-W. Liu. *No-go theorems for quantum resource purification: New approach and channel theory*. PRX Quantum, 3(1):010337 (2022)

Haar measure

We denote by

$$\mathbb{E}_U[f(U)] := \int_{U(d)} f(U) d\mu_H(U)$$

the expectation value of f over the Haar measure on $U(d) \ni U$, where μ_H is the **Haar measure** on the unitary group $U(d)$:

$$\int_{U(d)} f(U) d\mu_H(U) = \int_{U(d)} f(UV) d\mu_H(U) = \int_{U(d)} f(VU) d\mu_H(U) \quad V \in U(d)$$

Random purification channel

Let $\mathcal{H}_A \simeq \mathcal{H}_B$ be two isomorphic Hilbert spaces.

For any integer $n \geq 1$, let R_n be the operator on $(\mathcal{H}_A \otimes \mathcal{H}_B)^{\otimes n}$ defined as

$$R_n := \mathbb{E}_{U_B} \left[(\mathbb{1}_{A^n} \otimes U_B^{\otimes n}) \Gamma_{AB}^{\otimes n} (\mathbb{1}_{A^n} \otimes (U_B^\dagger)^{\otimes n}) \right].$$

The channel

$$\Lambda^{(n)}(\cdot) := \sqrt{R_n} (\cdot \otimes \mathbb{1}_{B^n}) \sqrt{R_n}$$

acts on i.i.d. states $\rho_A^{\otimes n}$, with $\rho_A \in \mathcal{H}_A$, as

$$\Lambda^{(n)}(\rho_A^{\otimes n}) = \mathbb{E}_{U_B} \left[(\mathbb{1}_{A^n} \otimes U_B^{\otimes n}) |\rho\rangle\langle\rho|_{AB}^{\otimes n} (\mathbb{1}_{A^n} \otimes U_B^{\dagger \otimes n}) \right],$$

where the expectation is taken over Haar-random unitaries U_B , and $|\rho\rangle\langle\rho|$ denotes any fixed purification of ρ .

Random purification channel

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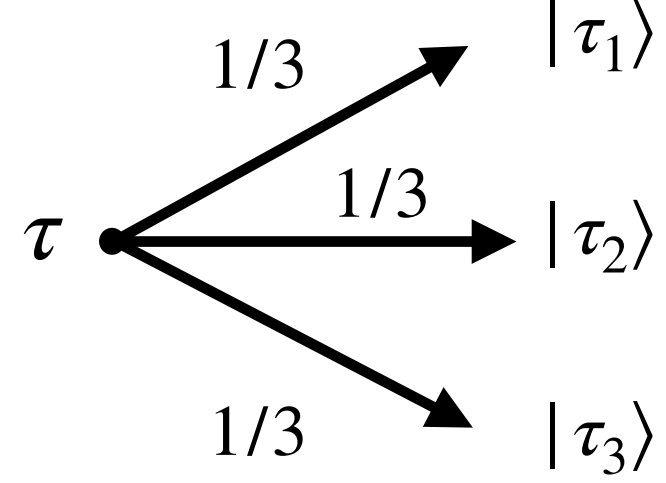
Random purification channel

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$(U_B^\dagger)^{\otimes n}$.

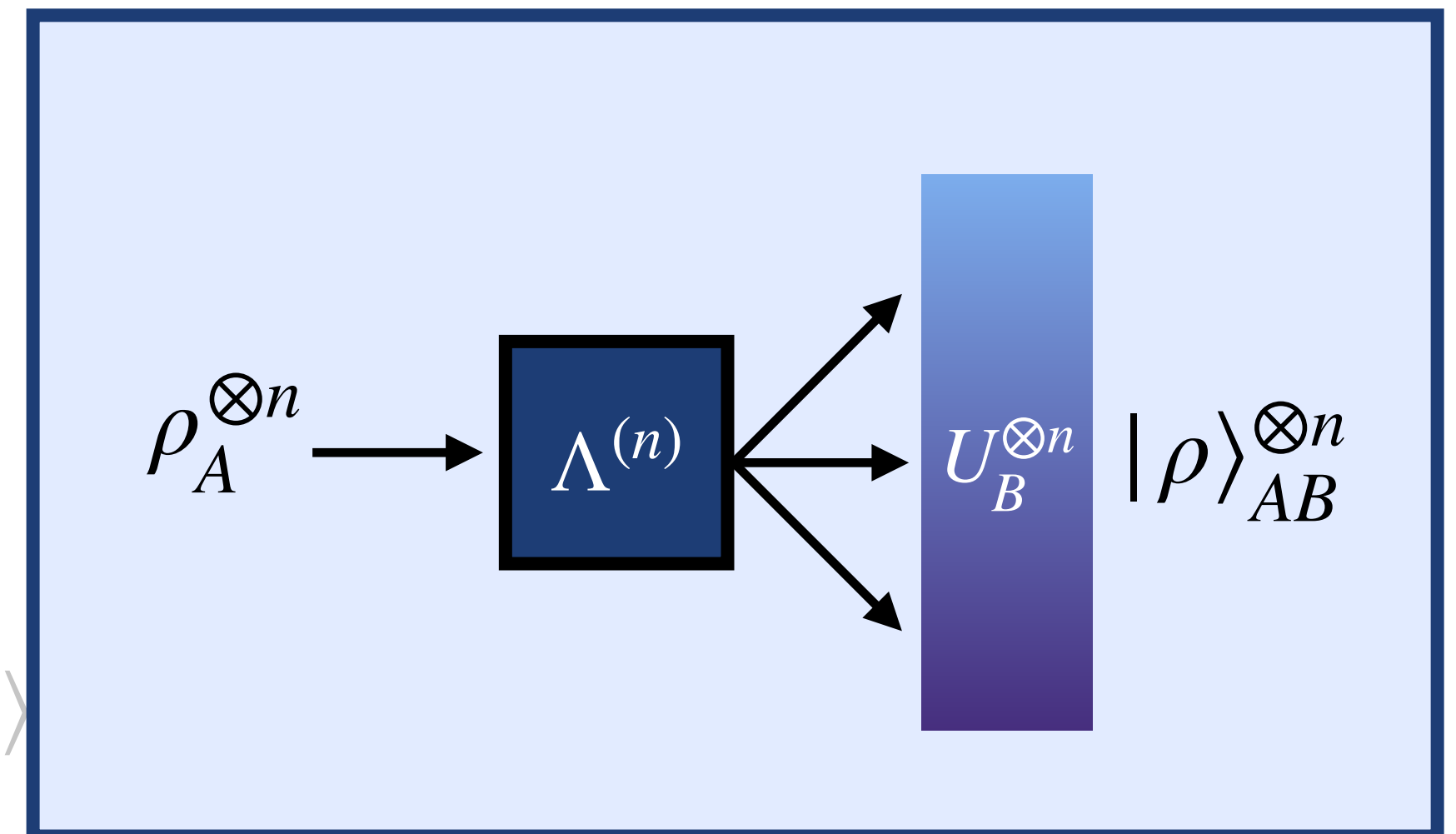
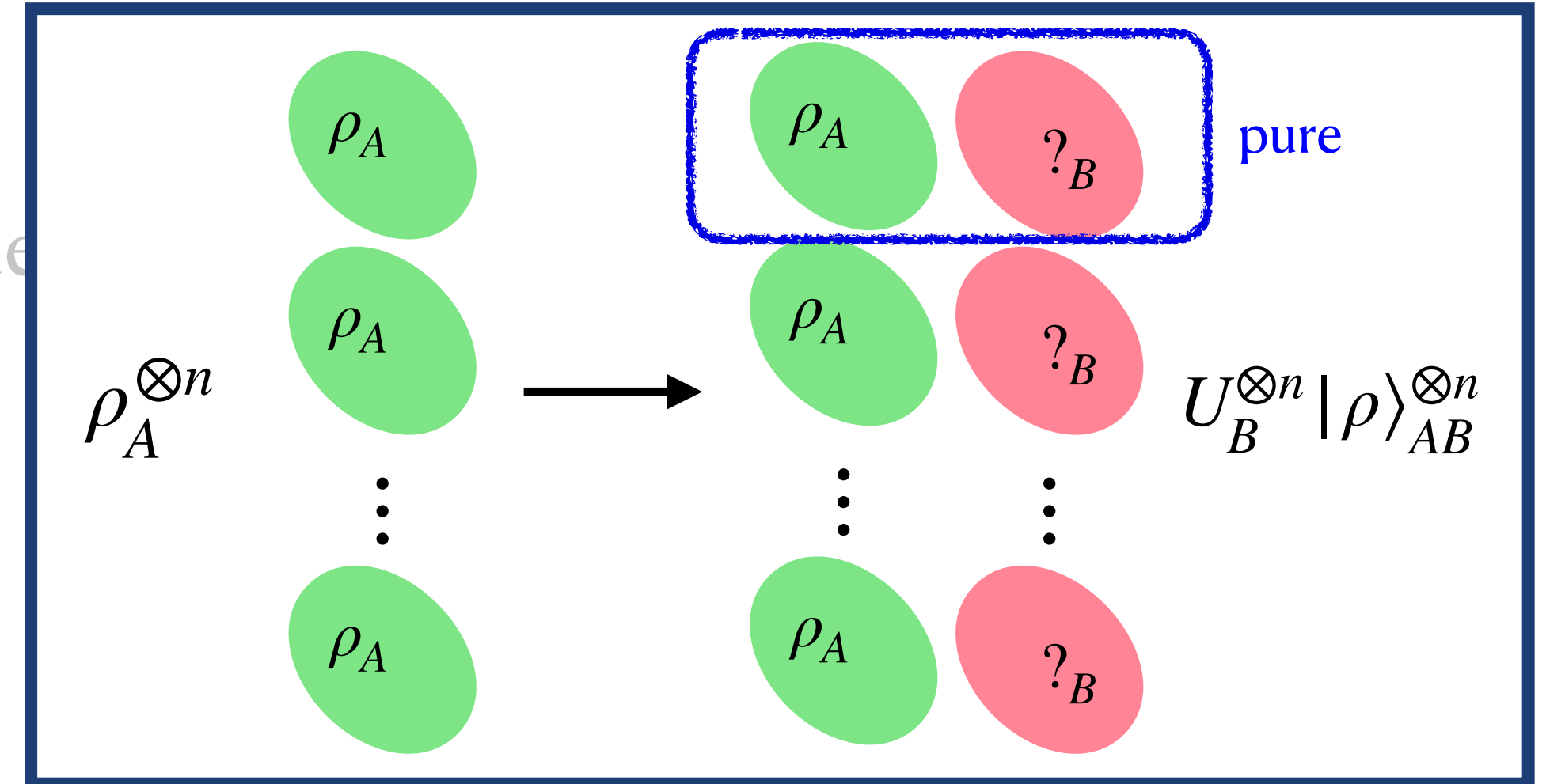
The

$$\tau = \frac{1}{3} |\tau_1\rangle\langle\tau_1| + \frac{1}{3} |\tau_2\rangle\langle\tau_2| + \frac{1}{3} |\tau_3\rangle\langle\tau_3|$$


acts

$$\Lambda^{(n)}(\rho_A^{\otimes n}) = \mathbb{E}_{U_B} \left[(\mathbb{I}_{A^n} \otimes U_B^{\otimes n}) |\rho\rangle\langle\rho|_{AB}^{\otimes n} (\mathbb{I}_{A^n} \otimes U_B^{\dagger \otimes n}) \right],$$

where the expectation is taken over Haar-random unitaries U_B , and $|\rho\rangle$



Simple proof

Let $\mathcal{H}_A \simeq \mathcal{H}_B$ be two isomorphic Hilbert spaces.

For any integer $n \geq 1$, let R_n be the operator on $(\mathcal{H}_A \otimes \mathcal{H}_B)^{\otimes n}$ defined as

$$R_n := \mathbb{E}_{U_B} \left[(\mathbb{1}_{A^n} \otimes U_B^{\otimes n}) \Gamma_{AB}^{\otimes n} (\mathbb{1}_{A^n} \otimes (U_B^\dagger)^{\otimes n}) \right].$$

$$\Gamma_{AB} := |\Gamma\rangle\langle\Gamma|_{AB}$$

$$|\Gamma\rangle_{AB} := \sum_{i=1}^d |e_i\rangle_A \otimes |e_i\rangle_B$$

Simple proof

Let $\mathcal{H}_A \simeq \mathcal{H}_B$ be two isomorphic Hilbert spaces.

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$$|\Gamma\rangle_{AB} := \sum_{i=1}^d |e_i\rangle_A \otimes |e_i\rangle_B$$

Lemma. Given $\rho_A \in \mathcal{D}(\mathcal{H}_A)$

$$[\rho_A^{\otimes n} \otimes \mathbb{1}_{B^n}, R_n] = 0$$

Standard purification. Given $\rho_A \in \mathcal{D}(\mathcal{H}_A)$

$$|\rho^{\text{std}}\rangle_{AB} := (\sqrt{\rho_A} \otimes \mathbb{1}_B) |\Gamma\rangle_{AB}$$

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The channel

$$\Lambda^{(n)}(\cdot) := \sqrt{R_n} (\cdot \otimes \mathbb{I}_{B^n}) \sqrt{R_n}$$

$$\begin{aligned} \Lambda^{(n)}(\rho_A^{\otimes n}) &= \sqrt{R_n} \rho_A^{\otimes n} \sqrt{R_n} = \sqrt{R_n} \sqrt{\rho_A^{\otimes n}} \sqrt{\rho_A^{\otimes n}} \sqrt{R_n} \\ &= \sqrt{\rho_A^{\otimes n}} \sqrt{R_n} \sqrt{R_n} \sqrt{\rho_A^{\otimes n}} = \sqrt{\rho_A^{\otimes n}} R_n \sqrt{\rho_A^{\otimes n}} \\ &= \mathbb{E}_{U_B} \left[(\mathbb{I}_{A^n} \otimes U_B^{\otimes n}) (\sqrt{\rho_A} \Gamma_{AB} \sqrt{\rho_A})^{\otimes n} (\mathbb{I}_{A^n} \otimes U_B^{\dagger \otimes n}) \right] \\ &= \mathbb{E}_{U_B} \left[(\mathbb{I}_{A^n} \otimes U_B^{\otimes n}) |\rho^{\text{std}}\rangle\langle\rho^{\text{std}}|_{AB}^{\otimes n} (\mathbb{I}_{A^n} \otimes U_B^{\dagger \otimes n}) \right]. \end{aligned}$$

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$$[\rho_A^{\otimes n} \otimes \mathbb{1}_{B^n}, R_n] = 0$$

STEP 1. $[R_n, U_A^{\otimes n} \otimes \mathbb{1}_{B^n}] = 0$

$$\begin{aligned} \text{Proof. } \left(U_A^{\otimes n} \otimes \mathbb{1}_{B^n} \right) R_n \left(U_A^{\otimes n} \otimes \mathbb{1}_{B^n} \right)^\dagger &= \mathbb{E}_{W_B} \left[\left((U_A \otimes W_B) \Gamma_{AB} (U_A \otimes W_B)^\dagger \right)^{\otimes n} \right] \\ &= \mathbb{E}_{W_B} \left[\left((\mathbb{1}_A \otimes W_B U_A^\dagger) \Gamma_{AB} (\mathbb{1}_A \otimes W_B U_A^\dagger)^\dagger \right)^{\otimes n} \right] = \mathbb{E}_{W'_B} \left[\left((\mathbb{1}_A \otimes W'_B) \Gamma_{AB} (\mathbb{1}_A \otimes W'_B)^\dagger \right)^{\otimes n} \right] = R_n, \end{aligned}$$

Simple proof

Let $\mathcal{H}_A \simeq \mathcal{H}_B$ be two isomorphic Hilbert spaces.

For any integer $n \geq 1$, let R_n be the operator on $(\mathcal{H}_A \otimes \mathcal{H}_B)^{\otimes n}$ defined as

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STEP 1. $[R_n, U_A^{\otimes n} \otimes \mathbb{1}_{B^n}] = 0$

STEP 2. $\left[R_n, \left(\sum_{j=1}^n H^{(j)} \right) \otimes \mathbb{1}_{B^n} \right] = 0$ where $H^{(j)} := \mathbb{1}^{\otimes(j-1)} \otimes H \otimes \mathbb{1}^{\otimes(n-j)}$, with H arbitrary Hermitian operator

Proof. $[R_n, (e^{itH})^{\otimes n} \otimes \mathbb{1}_{B^n}] = 0, \quad t \in \mathbb{R} \quad \rightarrow \quad \left. \frac{d}{dt} \right|_{t=0} \rightarrow \left[R_n, \left(\sum_{j=1}^n H^{(j)} \right) \otimes \mathbb{1}_{B^n} \right] = 0$

Simple proof

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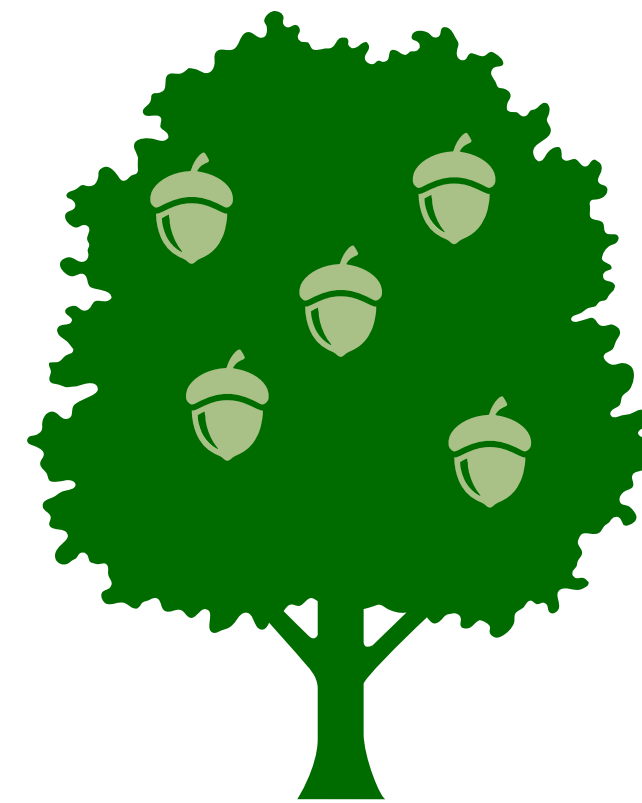
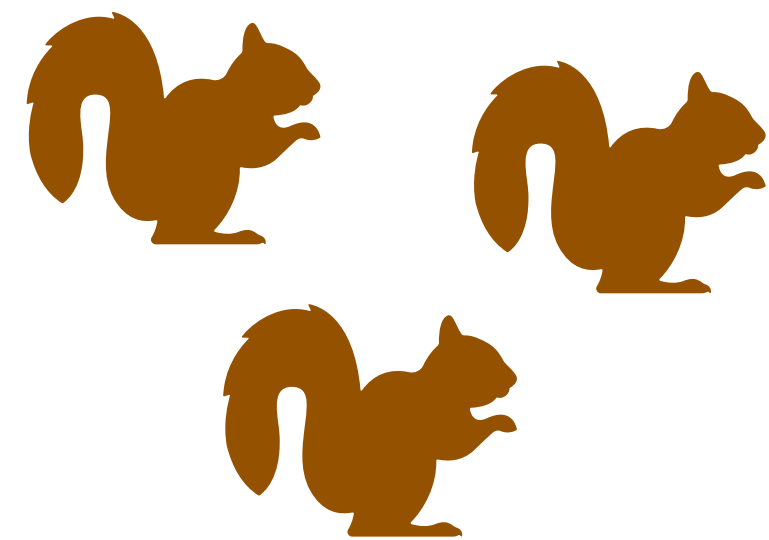
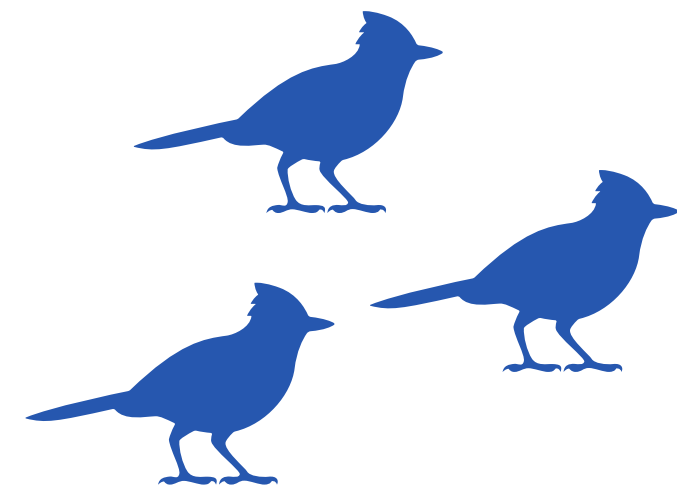
STEP 2. $\left[R_n, \left(\sum_{j=1}^n H^{(j)} \right) \otimes \mathbb{1}_{B^n} \right] = 0$ where $H^{(j)} := \mathbb{1}^{\otimes(j-1)} \otimes H \otimes \mathbb{1}^{\otimes(n-j)}$, with H arbitrary Hermitian operator

STEP 3. $\left[R_n, \rho_A^{\otimes n} \otimes \mathbb{1}_{B^n} \right] = 0$

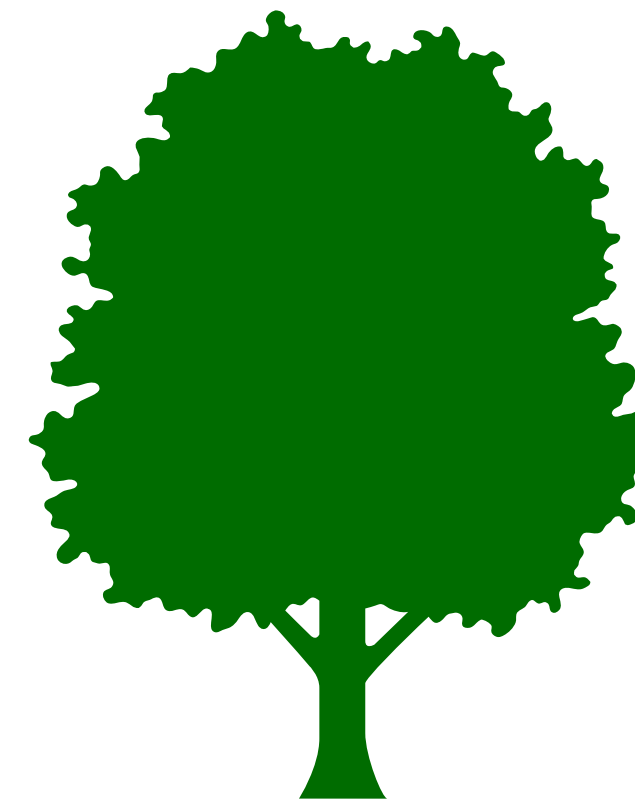
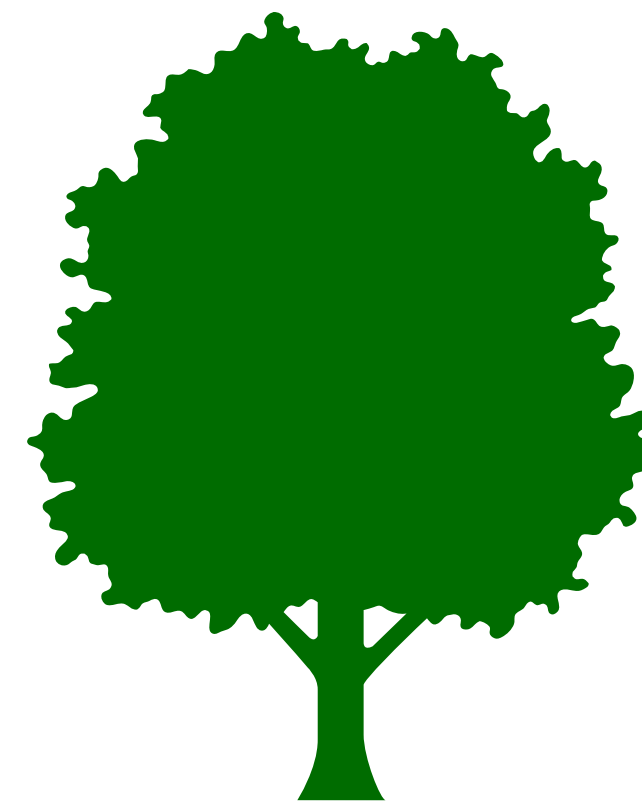
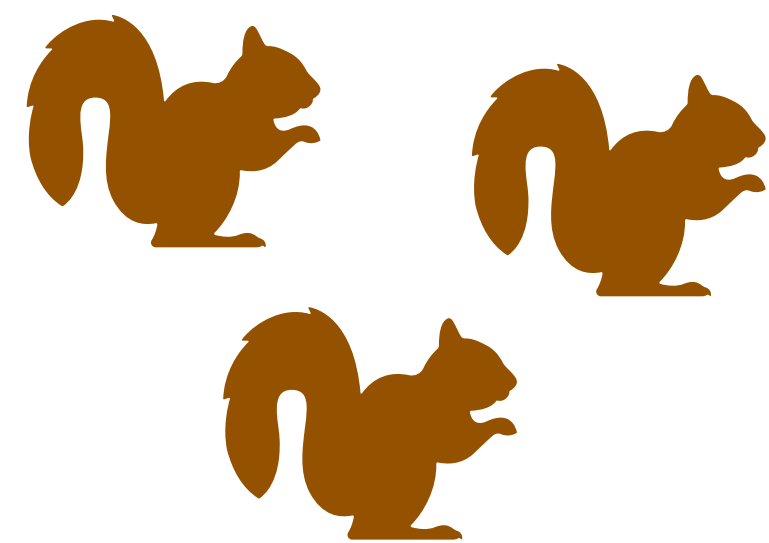
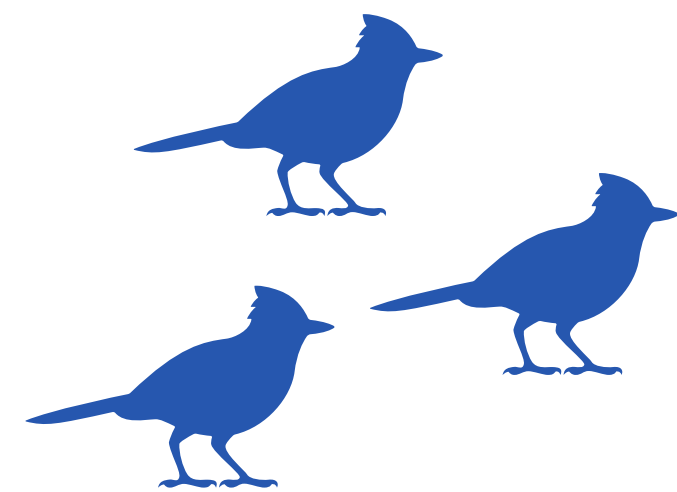
Proof. $0 = \left[R_n, \exp\left(\sum_{j=1}^n H^{(j)} \right) \otimes \mathbb{1}_{B^n} \right] = \left[R_n, (e^H)^{\otimes n} \otimes \mathbb{1}_{B^n} \right]$

$$\rho_A > 0 \quad \rightarrow \quad \rho_A = \frac{e^H}{\text{Tr } e^H} \quad \rightarrow \quad \left[R_n, \rho_A^{\otimes n} \otimes \mathbb{1}_{B^n} \right] = 0 \quad (\text{extension to } \rho_A \geq 0 \text{ by continuity})$$

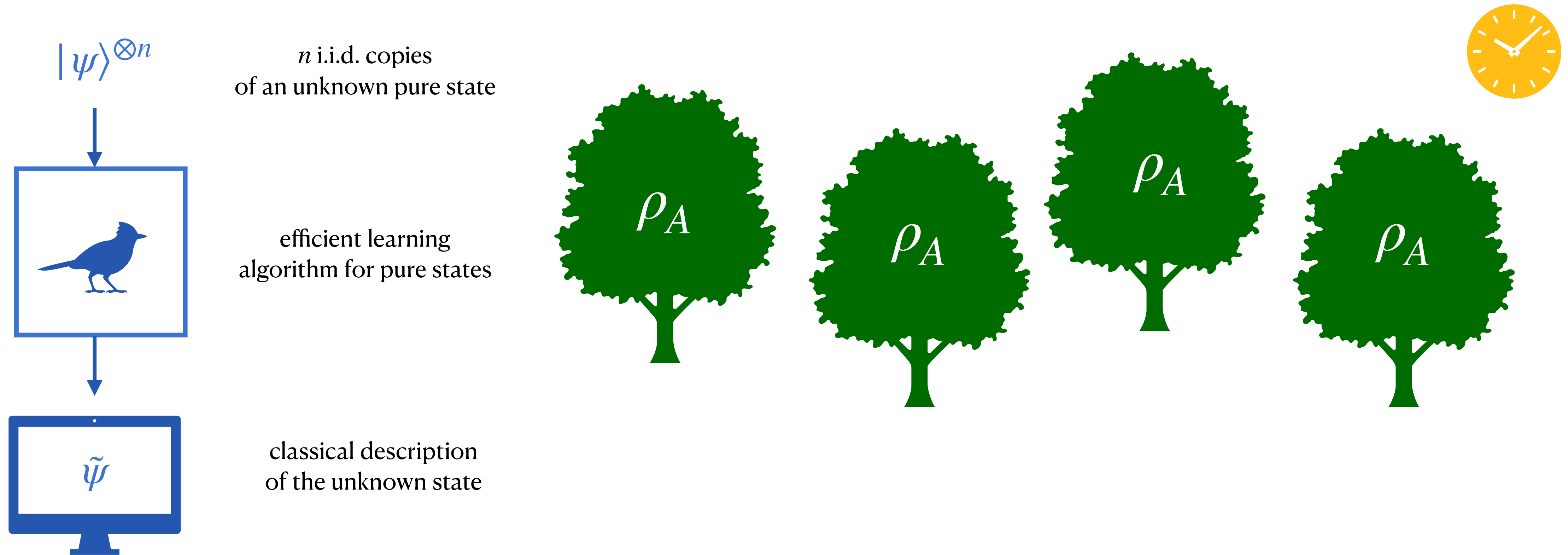
“Acorn trick”



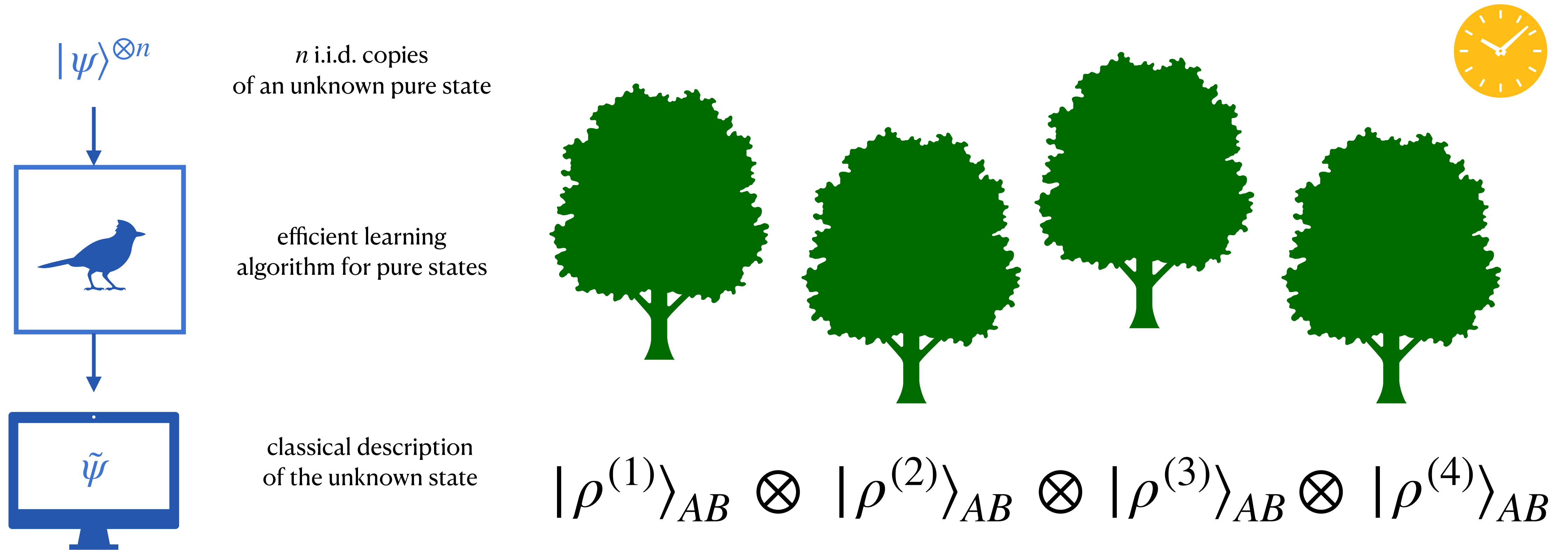
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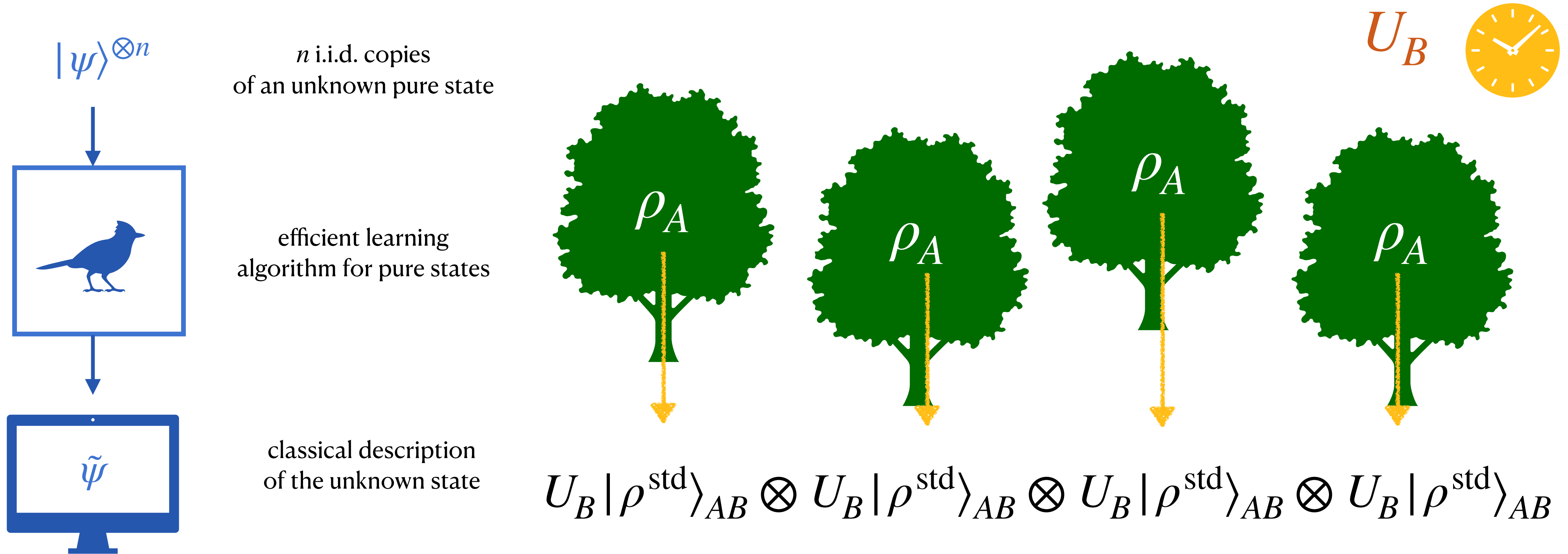
Applications in learning theory



Applications in learning theory

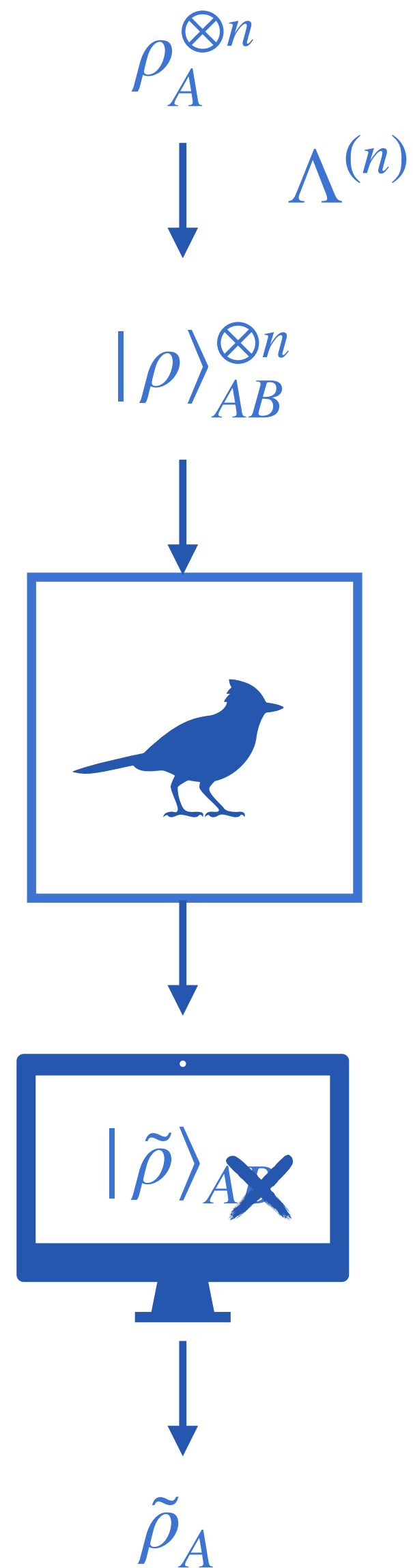


Applications in learning theory



$$\Lambda^{(n)}(\rho_A^{\otimes n}) = \mathbb{E}_{U_B} \left[(\mathbb{I}_{A^n} \otimes U_B^{\otimes n}) |\rho\rangle\langle\rho|_{AB}^{\otimes n} (\mathbb{I}_{A^n} \otimes U_B^{\dagger \otimes n}) \right]$$

Applications in learning theory



1. State tomography

A. Pelecanos, J. Spilecki, E. Tang, and J. Wright.

Mixed state tomography reduces to pure state tomography. arXiv:2511.15806 (2025)

$$N = O\left(\frac{rd + \log(1/\delta)}{\varepsilon}\right)$$

2. Channel tomography

A. A. Mele, L. Bittel

Optimal learning of quantum channels in diamond distance. arXiv:2512.10214 (2025)

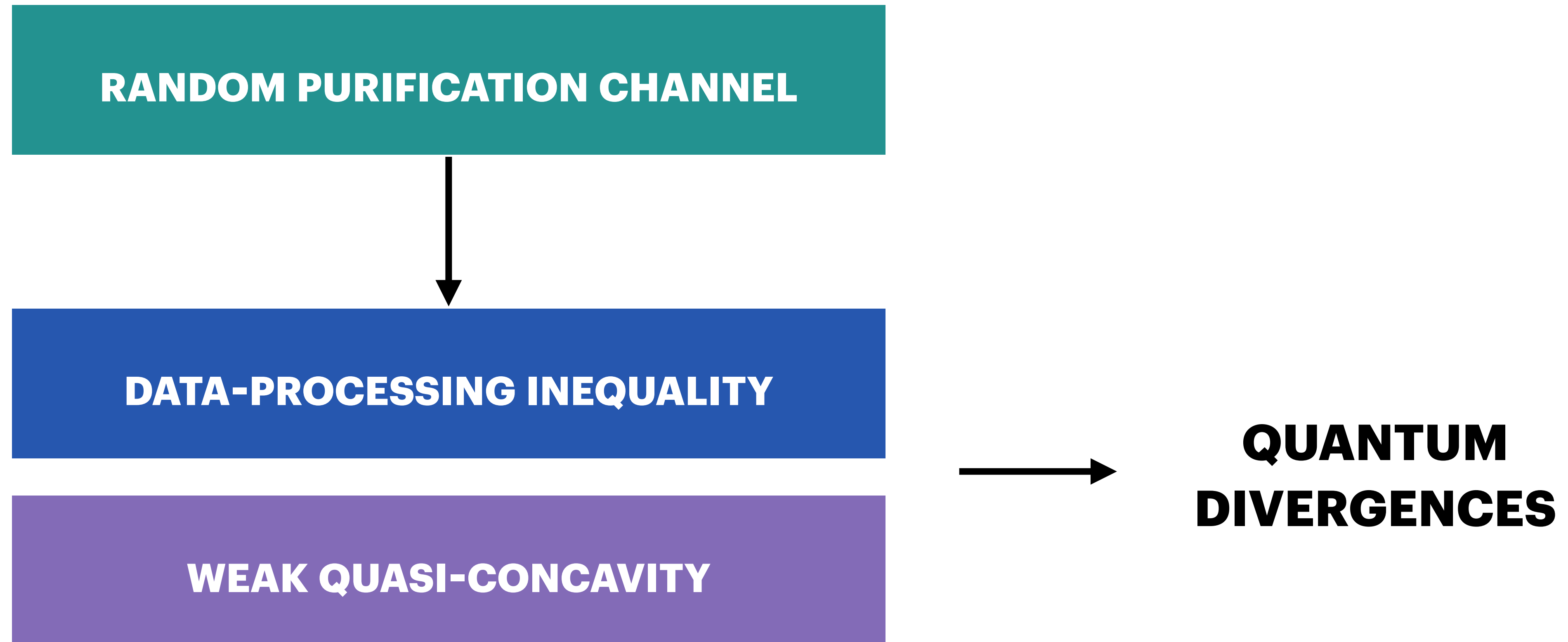
$$N = O\left(\frac{d_{\text{in}}d_{\text{out}}^r}{\varepsilon^2}\right)$$

3. Fidelity estimation

T. Utsumi, Y. Nakata, Q. Wang, and R. Takagi.

Quantum algorithms for Uhlmann transformation. arXiv:2509.03619 (2025)

Applications in quantum Shannon theory



Weak quasi-concavity

Quantum divergence

$$\mathbb{D}(\rho \parallel \sigma) \geq \mathbb{D}(\Lambda(\rho) \parallel \Lambda(\sigma)) \quad \Lambda \text{ quantum channel}$$

Weak quasi-concavity

$$\mathbb{D}\left(\sum_{i=1}^N p_i \rho_i \parallel \sigma\right)$$

Weak quasi-concavity

Quantum divergence

$$\mathbb{D}(\rho \parallel \sigma) \geq \mathbb{D}(\Lambda(\rho) \parallel \Lambda(\sigma)) \quad \Lambda \text{ quantum channel}$$

Weak quasi-concavity

$$\mathbb{D}\left(\sum_{i=1}^N p_i \rho_i \parallel \sigma\right) \geq \sum_{i=1}^N p_i D(\rho_i \parallel \sigma)$$

Weak quasi-concavity

Quantum divergence

$$\mathbb{D}(\rho \parallel \sigma) \geq \mathbb{D}(\Lambda(\rho) \parallel \Lambda(\sigma)) \quad \Lambda \text{ quantum channel}$$

Weak quasi-concavity

$$\mathbb{D}\left(\sum_{i=1}^N p_i \rho_i \parallel \sigma\right) \geq \min_{1 \leq i \leq N} \mathbb{D}(\rho_i \parallel \sigma) - \log \text{poly}_d(N, s_\sigma)$$

for any finite ensemble of states $\{(p_i, \rho_i)\}_{i=1, \dots, N}$ on $\mathcal{H}^{\otimes n}$ and any $\sigma \in \mathcal{D}(\mathcal{H}^{\otimes n})$

Weak quasi-concavity

Quantum divergence $\mathbb{D}(\rho \parallel \sigma) \geq \mathbb{D}(\Lambda(\rho) \parallel \Lambda(\sigma))$ Λ quantum channel

Weak quasi-concavity $\mathbb{D}\left(\sum_{i=1}^N p_i \rho_i \parallel \sigma\right) \geq \min_{1 \leq i \leq N} \mathbb{D}(\rho_i \parallel \sigma) - \log \text{poly}_d(N, s_\sigma)$

for any finite ensemble of states $\{(p_i, \rho_i)\}_{i=1, \dots, N}$ on $\mathcal{H}^{\otimes n}$ and any $\sigma \in \mathcal{D}(\mathcal{H}^{\otimes n})$

Quantum relative entropy $D(\rho \parallel \sigma) := \text{Tr}[\rho(\log \rho - \log \sigma)]$

Weak quasi-concavity $D\left(\sum_i p_i \rho_i \parallel \sigma\right) \geq \sum_i p_i D(\rho_i \parallel \sigma) + \sum_i p_i \log p_i$

- sandwiched Rényi divergences \tilde{D}_α of order $\alpha \in (0, \infty]$
- measured Rényi divergences $D_{M, \alpha}$ of order $\alpha \in (0, \infty]$
- #-Rényi divergences $D_\alpha^\#$ of order $\alpha \in (1, \infty]$

Applications in quantum Shannon theory

Uhlmann's theorem for quantum divergences

Fidelity: $F(\rho, \sigma) := \left\| \sqrt{\rho} \sqrt{\sigma} \right\|_1^2$

For pure states $F(|\psi\rangle, |\phi\rangle) = |\langle \psi | \phi \rangle|^2$

$$F(\rho, \sigma) = \max_{|\sigma\rangle} |\langle \rho | \sigma \rangle|^2$$

Uhlmann's theorem

$$\mathbb{D}^\infty(\rho_A \| \sigma_A) = \mathbb{D}^\infty(\rho_{AB} \| \mathcal{C}_{AB}^{\sigma_A})$$

Uhlmann's theorem for divergences

G. Mazzola, D. Sutter, and R. Renner, *Uhlmann's theorem for relative entropies*. IEEE Trans. Inf. Theory, 71(9):7039–7051 (2025)

K. Fang, H. Fawzi, and O. Fawzi, *Variational expressions and Uhlmann theorem for measured divergences*. IEEE Trans. Inf. Theory, 72(3):1751-1760 (2026)

Uhlmann's theorem for divergences

in one line

Let $\mathbb{D}(\cdot \| \cdot)$ be a divergence that obeys **weak quasi-concavity**.

Then, given ρ_A and σ_A in $\mathcal{D}(\mathcal{H}_A)$, for any arbitrary purification $\bar{\psi}_{AB}$ of ρ_A we have

$$\mathbb{D}^\infty(\rho_A \| \sigma_A) = \mathbb{D}^\infty(\bar{\psi}_{AB} \| \mathcal{C}_{AB}^{\sigma_A})$$

regularisation $\mathbb{D}^\infty(\rho_A \| \sigma_A) := \lim_{n \rightarrow \infty} \frac{1}{n} \mathbb{D}(\rho_A^{\otimes n} \| \sigma_A^{\otimes n})$

$$\mathbb{D}^\infty(\rho_A \| \sigma_A) := \lim_{n \rightarrow \infty} \frac{1}{n} \inf_{\sigma_{A^n B^n}} \mathbb{D}(\bar{\psi}_{AB}^{\otimes n} \| \sigma_{A^n B^n})$$

$$\uparrow$$

$$\text{Tr}_{B^n}[\sigma_{A^n B^n}] = \sigma_A^{\otimes n}$$

RANDOM PURIFICATION CHANNEL

$$\frac{1}{n} \mathbb{D}(\bar{\psi}^{\otimes n} \| \mathcal{C}_n^\sigma) \geq \frac{1}{n} \mathbb{D}(\rho_A^{\otimes n} \| \sigma_A^{\otimes n}) \geq \frac{1}{n} \mathbb{D}\left(\mathbb{E}_{|\psi_\rho\rangle} \psi_\rho^{\otimes n} \parallel \mathbb{E}_{|\phi_\sigma\rangle} \phi_\sigma^{\otimes n}\right) \gtrsim \frac{1}{n} \min_{|\psi_\rho\rangle} \mathbb{D}\left(\psi_\rho^{\otimes n} \parallel \mathbb{E}_{|\phi_\sigma\rangle} \phi_\sigma^{\otimes n}\right) = \frac{1}{n} \mathbb{D}\left(\bar{\psi}^{\otimes n} \parallel \mathbb{E}_{|\phi_\sigma\rangle} \phi_\sigma^{\otimes n}\right) \geq \frac{1}{n} \mathbb{D}(\bar{\psi}^{\otimes n} \| \mathcal{C}_n^\sigma),$$

DATA-PROCESSING INEQUALITY

WEAK QUASI-CONCAVITY

left invariance of Haar measure

Applications in quantum Shannon theory

RANDOM PURIFICATION CHANNEL



DATA-PROCESSING INEQUALITY

WEAK QUASI-CONCAVITY

**Are there any other
settings in which this
approach can be helpful?**



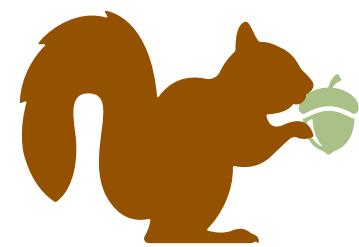
**QUANTUM
DIVERGENCES**

Random purification channel **BEYOND IID**

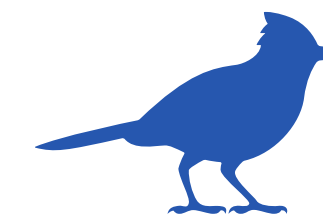
Permutation symmetry

$$\rho_{A^n} = U_\pi \rho_{A^n} U_\pi^\dagger \quad \forall \pi \in S_n$$

$$\Lambda^{(n)}(\rho_{A^n}) = \mathbb{E}_{U_B} \left[(\mathbb{I}_{A^n} \otimes U_B^{\otimes n}) |\rho\rangle\langle\rho|_{A^n B^n} (\mathbb{I}_{A^n} \otimes U_B^{\dagger \otimes n}) \right]$$



Thank you for your attention!



FG, F. A. Mele, L. Lami. *Random purification channel made simple*. arXiv:2511.23451 (2025)

F. A. Mele, FG, S. Chen, M. Fanizza, L. Lami. *Random purification channel for passive Gaussian bosons*. arXiv:2512.16878 (2025)

M. Walter, F. Witteveen. *A random purification channel for arbitrary symmetries with applications to fermions and bosons*. arXiv:2512.15690 (2025)